

U.S. DEPARTMENT OF COMMERCE  
Technology Administration  
National Institute of Standards and Technology



## FIPS PUB 180

### FEDERAL INFORMATION PROCESSING STANDARDS PUBLICATION

## SECURE HASH STANDARD

CATEGORY: COMPUTER SECURITY

1993 MAY 11

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Federal Information Processing Standards Publication

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### SECURE HASH STANDARD

Federal Information Processing Standard Publication 180, Secure Hash Standard, is published by the National Institute of Standards and Technology (NIST) under the authority granted to the Secretary of Commerce pursuant to Section 11(b)(d) of the Federal Property and Administrative Services Act of 1949 as amended by the Computer Security Act of 1987, Public Law 100-286.

#### Secure Hash Standard

##### CATEGORY: COMPUTER SECURITY

**Explanation:** This Standard specifies a Secure Hash Algorithm (SHA) for generating a condensed representation of a message or a data file. It takes any message less than  $2^{64}$  bits in input, the SHA produces a 160-bit output called a message digest. The message digest can then be input to some other process which can then verify the message digest against the signature for the message. The same hash algorithm is used in digital signatures as well as in the generation of checksums to verify the integrity of data files.

**Computer Systems Laboratory**  
National Institute of Standards and Technology  
Gaithersburg, MD 20899

The SHA is called secure because it is computationally infeasible to find a message which corresponds to a given message digest, or to find two messages which produce the same message digest. Any change in a message longer than 160 bits, with very high probability, results in a different message digest, and the digest can be used to verify the message. The SHA is based on principles similar to those used by Professor Ronald L. Rivest in developing the MD5 message digest algorithm, and is closely related to the SHA-256.

Issued May 11, 1993



**U.S. Department of Commerce**  
Ronald H. Brown, Secretary

National Institute of Standards and Technology  
Raymond G. Kammer, Acting Director

U.S. Government Printed Ogo 1993

W/McGraw-Hill Book Company

Standardized to sufficient length

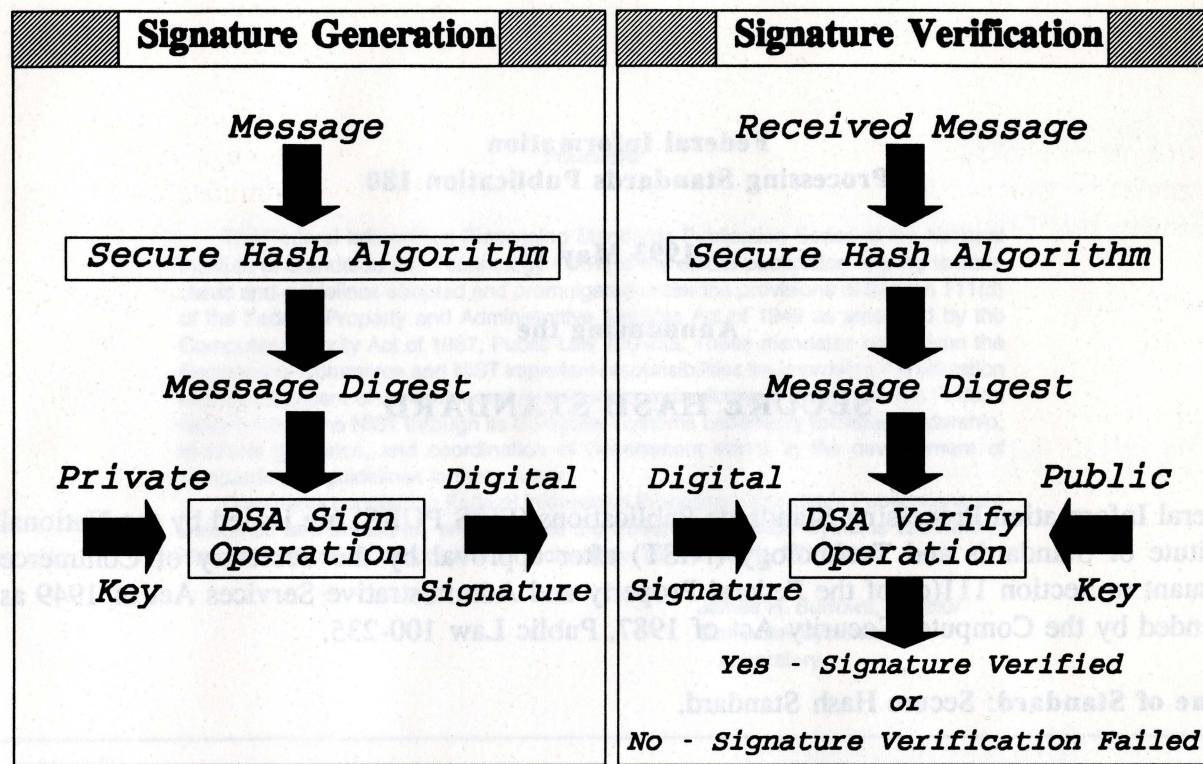
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(2001, 1st edn) 1993

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**Figure 1: Using the SHA with the DSA**

**Maintenance Agency:** U.S. Department of Commerce, National Institute of Standards and Technology, Computer Systems Laboratory.

**Applicability:** This standard is applicable to all Federal departments and agencies for the protection of unclassified information that is not subject to section 2315 of Title 10, United States Code, or section 3502(2) of Title 44, United States Code. This standard is required for use with the Digital Signature Algorithm (DSA) as specified in the Digital Signature Standard (DSS) and whenever a secure hash algorithm is required for federal applications. Private and commercial organizations are encouraged to adopt and use this standard.

**Applications:** The SHA may be used with the DSA in electronic mail, electronic funds transfer, software distribution, data storage, and other applications which require data integrity assurance and data origin authentication. The SHA may also be used whenever it is necessary to generate a condensed version of a message.

**Implementations:** The SHA may be implemented in software, firmware, hardware, or any combination thereof. Only implementations of the SHA that are validated by NIST will be

considered as complying with this standard. Information about the requirements for validating implementations of this standard can be obtained from the National Institute of Standards and Technology, Computer Systems Laboratory, Attn: SHS Validation, Gaithersburg, MD 20899.

**Export Control:** Implementations of this standard are subject to Federal Government export controls as specified in Title 15, Code of Federal Regulations, Parts 768 through 799. Exporters are advised to contact the Department of Commerce, Bureau of Export Administration for more information.

**Patents:** Implementations of the SHA in this standard may be covered by U.S. and foreign patents.

**Implementation Schedule:** This standard becomes effective October 15, 1993.

**Specifications:** Federal Information Processing Standard (FIPS 180) Secure Hash Standard (affixed).

**Cross Index:**

- a. FIPS PUB 46-1, Data Encryption Standard.
- b. FIPS PUB 73, Guidelines for Security of Computer Applications.
- c. Draft FIPS PUB 140-1, Security Requirements for Cryptographic Modules.
- d. FIPS PUB XX, Digital Signature Standard.

**Qualifications:** While it is the intent of this standard to specify a secure hash algorithm, conformance to this standard does not assure that a particular implementation is secure. The responsible authority in each agency or department shall assure that an overall implementation provides an acceptable level of security. This standard will be reviewed every five years in order to assess its adequacy.

**Waiver Procedure:** Under certain exceptional circumstances, the heads of Federal departments and agencies may approve waivers to Federal Information Processing Standards (FIPS). The head of such agency may redelegate such authority only to a senior official designated pursuant to section 3506(b) of Title 44, United States Code. Waiver shall be granted only when:

- a. Compliance with a standard would adversely affect the accomplishment of the mission of an operator of a Federal computer system; or
- b. Compliance with a standard would cause a major adverse financial impact on the operator which is not offset by Government-wide savings.

Agency heads may act upon a written waiver request containing the information detailed above. Agency heads may also act without a written waiver request when they determine that conditions for meeting the standard cannot be met. Agency heads may approve waivers only by a written decision which explains the basis on which the agency head made the required finding(s). A copy of each decision, with procurement sensitive or classified portions clearly identified, shall be sent to: National Institute of Standards and Technology; ATTN: FIPS Waiver Decisions, Technology Building, Room B-154, Gaithersburg, MD 20899.

In addition, notice of each waiver granted and each delegation of authority to approve waivers shall be sent promptly to the Committee on Government Operations of the House of Representatives and the Committee on Government Affairs of the Senate and shall be published promptly in the Federal Register.

When the determination on a waiver applies to the procurement of equipment and/or services, a notice of the waiver determination must be published in the Commerce Business Daily as a part of the notice of solicitation for offers of an acquisition or, if the waiver determination is made after that notice is published, by amendment to such notice.

A copy of the waiver, any supporting documents, the document approving the waiver and any accompanying documents, with such deletions as the agency is authorized and decides to make under 5 United States Code Section 552(b), shall be part of the procurement documentation and retained by the agency.

**Where to Obtain Copies of the Standard:** Copies of this publication are for sale by the National Technical Information Service, U.S. Department of Commerce, Springfield, VA 22161. When ordering, refer to Federal Information Processing Standards Publication 180 (FIPSPUB180), and identify the title. When microfiche is desired, this should be specified. Prices are published by NTIS in current catalogs and other issuances. Payment may be made by check, money order, deposit account or charged to a credit card accepted by NTIS.

Implementation. The SIA may implement FIPS 180 through review or any combination thereof. Only implementations of the SIA that are validated by NIST will be

**Federal Information****Processing Standards Publication 180**

1993 May 11

**Specifications for the****SECURE HASH STANDARD****1. INTRODUCTION**

The Secure Hash Algorithm (SHA) is required for use with the Digital Signature Algorithm (DSA) as specified in the Digital Signature Standard (DSS) and whenever a secure hash algorithm is required for federal applications. For a message of length  $< 2^{64}$  bits, the SHA produces a 160-bit condensed representation of the message called a message digest. The message digest is used during generation of a signature for the message. The SHA is also used to compute a message digest for the received version of the message during the process of verifying the signature. Any change to the message in transit will, with very high probability, result in a different message digest, and the signature will fail to verify.

The SHA is designed to have the following properties: it is computationally infeasible to find a message which corresponds to a given message digest, or to find two different messages which produce the same message digest.

**2. BIT STRINGS AND INTEGERS**

The following terminology related to bit strings and integers will be used:

- A *hex digit* is an element of the set {0, 1, ..., 9, A, ..., F}. A hex digit is the representation of a 4-bit string. **Example:** 7 = 0111, A = 1010.
- A *word* equals a 32-bit string which may be represented as a sequence of 8 hex digits. To convert a word to 8 hex digits each 4-bit string is converted to its hex equivalent as described in (a) above. **Example:**

1010 0001 0000 0011 1111 1110 0010 0011. = A103FE23

- An *integer between 0 and  $2^{32} - 1$  inclusive* may be represented as a word. The least significant four bits of the integer are represented by the right-most hex digit of the word representation. **Example:** the integer  $291 = 2^8 + 2^5 + 2^1 + 2^0 = 256 + 32 + 2 + 1$  is represented by the hex word, 00000123.

If  $z$  is an integer,  $0 \leq z < 2^{64}$ , then  $z = 2^{32}x + y$  where  $0 \leq x < 2^{32}$  and  $0 \leq y < 2^{32}$ . Since  $x$  and  $y$  can be represented as words  $X$  and  $Y$ , respectively,  $z$  can be represented as the pair of words  $(X, Y)$ .

- d. *block* = 512-bit string. A block (e.g.,  $B$ ) may be represented as a sequence of 16 words.

### 3. OPERATIONS ON WORDS

The following logical operators will be applied to words:

- a. *Bitwise logical word operations*

$$X \wedge Y = \text{bitwise logical "and" of } X \text{ and } Y.$$

$$X \vee Y = \text{bitwise logical "inclusive-or" of } X \text{ and } Y.$$

$$X \text{ XOR } Y = \text{bitwise logical "exclusive-or" of } X \text{ and } Y.$$

$$\sim X = \text{bitwise logical "complement" of } X.$$

**Example:**

$$\begin{array}{rcl} & 01101100101110011101001001111011 \\ \text{XOR} & 01100101110000010110100110110111 \\ \hline & 00001001011110001011101111001100 \end{array}$$

- b. The *operation*  $X + Y$  is defined as follows: words  $X$  and  $Y$  represent integers  $x$  and  $y$ , where  $0 \leq x < 2^{32}$  and  $0 \leq y < 2^{32}$ . For positive integers  $n$  and  $m$ , let  $n \bmod m$  be the remainder upon dividing  $n$  by  $m$ . Compute

$$z = (x + y) \bmod 2^{32}.$$

Then  $0 \leq z < 2^{32}$ . Convert  $z$  to a word,  $Z$ , and define  $Z = X + Y$ .

- c. The *circular left shift operation*  $S^n(X)$ , where  $X$  is a word and  $n$  is an integer with  $0 \leq n < 32$ , is defined by

$$S^n(X) = (X \ll n) \vee (X \gg 32-n).$$

In the above,  $X \ll n$  is obtained as follows: discard the left-most  $n$  bits of  $X$  and then pad the result with  $n$  zeroes on the right (the result will still be 32 bits).  $X \gg n$  is obtained by discarding the right-most  $n$  bits of  $X$  and then padding the result with  $n$

zeroes on the left. Thus  $S^n(X)$  is equivalent to a circular shift of  $X$  by  $n$  positions to the left.

#### 4. MESSAGE PADDING

The SHA is used to compute a message digest for a message or data file that is provided as input. The message or data file should be considered to be a bit string. The length of the message is the number of bits in the message (the empty message has length 0). If the number of bits in a message is a multiple of 8, for compactness we can represent the message in hex. The purpose of message padding is to make the total length of a padded message a multiple of 512. The SHA sequentially processes blocks of 512 bits when computing the message digest. The following specifies how this padding shall be performed. As a summary, a "1" followed by  $m$  "0"s followed by a 64-bit integer are appended to the end of the message to produce a padded message of length  $512 \times n$ . The 64-bit integer is  $l$ , the length of the original message. The padded message is then processed by the SHA as  $n$  512-bit blocks.

Suppose a message has length  $l < 2^{64}$ . Before it is input to the SHA, the message is padded on the right as follows:

- "1" is appended. **Example:** if the original message is "01010000", this is padded to "010100001".
- "0"s are appended. The number of "0"s will depend on the original length of the message. The last 64 bits of the last 512-bit block are reserved for the length  $l$  of the original message.

**Example:** Suppose the original message is the bit string

01100001 01100010 01100011 01100100 01100101.

After step (a) this gives

01100001 01100010 01100011 01100100 01100101 1.

Since  $l = 40$ , the number of bits in the above is 41 and 407 "0"s are appended, making the total now 448. This gives (in hex)

61626364	65800000	00000000	00000000
00000000	00000000	00000000	00000000
00000000	00000000	00000000	00000000
00000000	00000000		

- Obtain the 2-word representation of  $l$ , the number of bits in the original message. If  $l < 2^{32}$  then the first word is all zeroes. Append these two words to the padded message.

**Example:** Suppose the original message is as in (b). Then  $l = 40$  (note that  $l$  is computed before any padding). The two-word representation of 40 is hex 00000000 00000028. Hence the final padded message is hex

61626364	65800000	00000000	00000000
00000000	00000000	00000000	00000000
00000000	00000000	00000000	00000000
00000000	00000000	00000000	00000028.

The padded message will contain  $16n$  words for some  $n > 0$ . The padded message is regarded as a sequence of  $n$  blocks  $M_1, M_2, \dots, M_n$ , where each  $M_i$  contains 16 words and  $M_1$  contains the first characters (or bits) of the message.

## 5. FUNCTIONS USED

A sequence of logical functions  $f_0, f_1, \dots, f_{79}$  is used in the SHA. Each  $f_t$ ,  $0 \leq t \leq 79$ , operates on three 32-bit words and produces a 32-bit word as output.  $f_t$  is defined as follows: for words,  $B, C, D$ ,

$$f_t(B, C, D) = (B \wedge C) \vee (\neg B \wedge D) \quad (0 \leq t \leq 19)$$

$$f_t(B, C, D) = B \text{ XOR } C \text{ XOR } D \quad (20 \leq t \leq 39)$$

$$f_t(B, C, D) = (B \wedge C) \vee (B \wedge D) \vee (C \wedge D) \quad (40 \leq t \leq 59)$$

$$f_t(B, C, D) = B \text{ XOR } C \text{ XOR } D \quad (60 \leq t \leq 79).$$

## 6. CONSTANTS USED

A sequence of constant words  $K_0, K_1, \dots, K_{79}$  is used in the SHA. In hex these are given by

$$K_t = 5A827999 \quad (0 \leq t \leq 19)$$

$$K_t = 6ED9EBA1 \quad (20 \leq t \leq 39)$$

$$K_t = 8F1BBCDC \quad (40 \leq t \leq 59)$$

$$K_t = CA62C1D6 \quad (60 \leq t \leq 79).$$

## APPENDIX A: A SAMPLE MESSAGE AND ITS MESSAGE DIGEST

### 7. COMPUTING THE MESSAGE DIGEST

The message digest is computed using the final padded message. The computation uses two buffers, each consisting of five 32-bit words, and a sequence of eighty 32-bit words. The words of the first 5-word buffer are labeled A,B,C,D,E. The words of the second 5-word buffer are labeled  $H_0, H_1, H_2, H_3, H_4$ . The words of the 80-word sequence are labeled  $W_0, W_1, \dots, W_{79}$ . A single word buffer TEMP is also employed.

To generate the message digest, the 16-word blocks  $M_1, M_2, \dots, M_n$  defined in Section 4 are processed in order. The processing of each  $M_i$  involves 80 steps.

Before processing any blocks, the  $\{H_j\}$  are initialized as follows: in hex,

$$H_0 = 67452301$$

$$H_1 = EFCDAB89$$

$$H_2 = 98BADCFE$$

$$H_3 = 10325476$$

$$H_4 = C3D2E1F0.$$

Now  $M_1, M_2, \dots, M_n$  are processed. To process  $M_i$ , we proceed as follows:

- a. Divide  $M_i$  into 16 words  $W_0, W_1, \dots, W_{15}$ , where  $W_0$  is the left-most word.
- b. For  $t = 16$  to 79 let  $W_t = W_{t-3} \text{ XOR } W_{t-8} \text{ XOR } W_{t-14} \text{ XOR } W_{t-16}$ .
- c. Let  $A = H_0, B = H_1, C = H_2, D = H_3, E = H_4$ .
- d. For  $t = 0$  to 79 do
  - $TEMP = S^5(A) + f_t(B, C, D) + E + W_t + K_t;$
  - $E = D; D = C; C = S^{30}(B); B = A; A = TEMP;$
- e. Let  $H_0 = H_0 + A, H_1 = H_1 + B, H_2 = H_2 + C, H_3 = H_3 + D, H_4 = H_4 + E$ .

After processing  $M_n$ , the message digest is the 160-bit string represented by the 5 words

$$H_0 \ H_1 \ H_2 \ H_3 \ H_4.$$

## 8. ALTERNATE METHOD OF COMPUTATION

The above assumes that the sequence  $W_0, \dots, W_{79}$  is implemented as an array of eighty 32-bit words. This is efficient from the standpoint of minimization of execution time, since the addresses of  $W_{t-3}, \dots, W_{t-16}$  in step (b) are easily computed. If space is at a premium, an alternative is to regard  $\{W_t\}$  as a circular queue, which may be implemented using an array of sixteen 32-bit words  $W[0], \dots, W[15]$ . In this case, in hex let MASK = 0000000F. Then processing of  $M_i$  is as follows:

- a. Divide  $M_i$  into 16 words  $W[0], \dots, W[15]$ , where  $W[0]$  is the left-most word.
- b. Let  $A = H_0, B = H_1, C = H_2, D = H_3, E = H_4$ .
- c. For  $t = 0$  to 79 do
  - $s = t \wedge \text{MASK};$
  - if ( $t \geq 16$ )  $W[s] = W[(s + 13) \wedge \text{MASK}] \text{ XOR } W[(s + 8) \wedge \text{MASK}] \text{ XOR } W[(s + 2) \wedge \text{MASK}] \text{ XOR } W[s];$
  - $\text{TEMP} = S^5(A) + f_t(B, C, D) + E + W[s] + K_t;$
  - $E = D; D = C; C = S^{30}(B); B = A; A = \text{TEMP};$
- d. Let  $H_0 = H_0 + A, H_1 = H_1 + B, H_2 = H_2 + C, H_3 = H_3 + D, H_4 = H_4 + E$ .

## 9. COMPARISON OF METHODS

The methods of Sections 7 and 8 yield the same message digest. Although using the method of Section 8 saves sixty-four 32-bit words of storage, it is likely to lengthen execution time due to the increased complexity of the address computations for the  $\{W[t]\}$  in step (c). Other computation methods which give identical results may be implemented in conformance with the standard.

## APPENDIX A. A SAMPLE MESSAGE AND ITS MESSAGE DIGEST

This appendix is for informational purposes only and is not required to meet the standard.

Let the message be the ASCII binary-coded form of "abc", i.e.,

01100001 01100010 01100011.

This message has length  $l = 24$ . In step (a) of Section 4, we append "1". In step (b) we append 423 "0"s. In step (c) we append hex 00000000 00000018, the 2-word representation of 24. Thus the final padded message consists of one block, so that  $n = 1$  in the notation of Section 4. The single block has hex words

$W[0]$	=	61626380
$W[1]$	=	00000000
$W[2]$	=	00000000
$W[3]$	=	00000000
$W[4]$	=	00000000
$W[5]$	=	00000000
$W[6]$	=	00000000
$W[7]$	=	00000000
$W[8]$	=	00000000
$W[9]$	=	00000000
$W[10]$	=	00000000
$W[11]$	=	00000000
$W[12]$	=	00000000
$W[13]$	=	00000000
$W[14]$	=	00000000
$W[15]$	=	00000018.

The initial hex values of  $\{H_i\}$  are

$$H_0 = 67452301$$

$$H_1 = EFCDAB89$$

$$H_2 = 98BADCFE$$

$$H_3 = 10325476$$

$$H_4 = C3D2E1F0.$$

The hex values of A,B,C,D,E after pass t of the "for  $t = 0$  to 79" loop (step (d) of Section 7 or (c) of Section 8) are

	A	B	C	D	E
t = 0:	0116FC33	67452301	7BF36AE2	98BADC <sup>FE</sup>	10325476
t = 1:	8990536D	0116FC33	59D148C0	7BF36AE2	98BADC <sup>FE</sup>
t = 2:	A1390F08	8990536D	C045BF0C	59D148C0	7BF36AE2
t = 3:	CDD8E11B	A1390F08	626414DB	C045BF0C	59D148C0
t = 4:	CFD499DE	CDD8E11B	284E43C2	626414DB	C045BF0C
t = 5:	3FC7CA40	CFD499DE	F3763846	284E43C2	626414DB
t = 6:	993E30C1	3FC7CA40	B3F52677	F3763846	284E43C2
t = 7:	9E8C07D4	993E30C1	0FF1F290	B3F52677	F3763846
t = 8:	4B6AE328	9E8C07D4	664F8C30	0FF1F290	B3F52677
t = 9:	8351F929	4B6AE328	27A301F5	664F8C30	0FF1F290
t = 10:	FBDA9E89	8351F929	12DAB8CA	27A301F5	664F8C30
t = 11:	63188FE4	FBDA9E89	60D47E4A	12DAB8CA	27A301F5
t = 12:	4607B664	63188FE4	7EF6A7A2	60D47E4A	12DAB8CA
t = 13:	9128F695	4607B664	18C623F9	7EF6A7A2	60D47E4A
t = 14:	196BEE77	9128F695	1181ED99	18C623F9	7EF6A7A2
t = 15:	20BDD62F	196BEE77	644A3DA5	1181ED99	18C623F9
t = 16:	ED2FF4A3	20BDD62F	C65AFB9D	644A3DA5	1181ED99
t = 17:	565DF73C	ED2FF4A3	C82F758B	C65AFB9D	644A3DA5
t = 18:	550B1E7F	565DF73C	FB4BFD28	C82F758B	C65AFB9D
t = 19:	FE0F9E4B	550B1E7F	15977DCF	FB4BFD28	C82F758B
t = 20:	B4D4C943	FE0F9E4B	D542C79F	15977DCF	FB4BFD28
t = 21:	43993572	B4D4C943	FF83E792	D542C79F	15977DCF
t = 22:	F7106486	43993572	ED353250	FF83E792	D542C79F
t = 23:	775924E6	F7106486	90E64D5C	ED353250	FF83E792
t = 24:	45A7EF23	775924E6	BDC41921	90E64D5C	ED353250
t = 25:	CCEAD674	45A7EF23	9DD64939	BDC41921	90E64D5C
t = 26:	02D0C6D1	CCEAD674	D169FBC8	9DD64939	BDC41921
t = 27:	070C437F	02D0C6D1	333AB59D	D169FBC8	9DD64939
t = 28:	301E90BE	070C437F	40B431B4	333AB59D	D169FBC8
t = 29:	B898C685	301E90BE	C1C310DF	40B431B4	333AB59D
t = 30:	669723E2	B898C685	8C07A42F	C1C310DF	40B431B4
t = 31:	D9316F96	669723E2	6E2631A1	8C07A42F	C1C310DF
t = 32:	DB81A5C7	D9316F96	99A5C8F8	6E2631A1	8C07A42F
t = 33:	99C8DFB2	DB81A5C7	B64C5BE5	99A5C8F8	6E2631A1
t = 34:	6BE6AE07	99C8DFB2	F6E06971	B64C5BE5	99A5C8F8
t = 35:	C01CC62C	6BE6AE07	A67237EC	F6E06971	B64C5BE5
t = 36:	6433FDD0	C01CC62C	DAF9AB81	A67237EC	F6E06971
t = 37:	0A33CCF7	6433FDD0	3007318B	DAF9AB81	A67237EC
t = 38:	4BF58DC8	0A33CCF7	190CFF74	3007318B	DAF9AB81
t = 39:	EBBD5233	4BF58DC8	C28CF33D	190CFF74	3007318B
t = 40:	825A3460	EBBD5233	12FD6372	C28CF33D	190CFF74
t = 41:	B62CBB93	825A3460	FAEF548C	12FD6372	C28CF33D
t = 42:	AA3F9707	B62CBB93	20968D18	FAEF548C	12FD6372
t = 43:	FE1D0273	AA3F9707	ED8B2EE4	20968D18	FAEF548C
t = 44:	57AD526B	FE1D0273	EA8FE5C1	ED8B2EE4	20968D18
t = 45:	93EBBE3F	57AD526B	FF87409C	EA8FE5C1	ED8B2EE4
t = 46:	F9ADF47B	93EBBE3F	D5EB549A	FF87409C	EA8FE5C1

t = 47:	875586D2	F9ADF47B	E4FAEF8F	D5EB549A	FF87409C
t = 48:	D0A22FFB	875586D2	FE6B7D1E	E4FAEF8F	D5EB549A
t = 49:	C12B6426	D0A22FFB	A1D561B4	FE6B7D1E	E4FAEF8F
t = 50:	EBC90281	C12B6426	F4288BFE	A1D561B4	FE6B7D1E
t = 51:	E7D0EC05	EBC90281	B04AD909	F4288BFE	A1D561B4
t = 52:	7CB98E55	E7D0EC05	7AF240A0	B04AD909	F4288BFE
t = 53:	0D48DBA2	7CB98E55	79F43B01	7AF240A0	B04AD909
t = 54:	C2D477BF	0D48DBA2	5F2E6395	79F43B01	7AF240A0
t = 55:	236BD48D	C2D477BF	835236E8	5F2E6395	79F43B01
t = 56:	9B4364D6	236BD48D	F0B51DEF	835236E8	5F2E6395
t = 57:	5B8C33C9	9B4364D6	48DAF523	F0B51DEF	835236E8
t = 58:	BE2A4656	5B8C33C9	A6D0D935	48DAF523	F0B51DEF
t = 59:	8FF296DB	BE2A4656	56E30CF2	A6D0D935	48DAF523
t = 60:	C10C8993	8FF296DB	AF8A9195	56E30CF2	A6D0D935
t = 61:	6AC23CBF	C10C8993	E3FCA5B6	AF8A9195	56E30CF2
t = 62:	0708247D	6AC23CBF	F0432264	E3FCA5B6	AF8A9195
t = 63:	35D201F8	0708247D	DAB08F2F	F0432264	E3FCA5B6
t = 64:	969B2FC8	35D201F8	41C2091F	DAB08F2F	F0432264
t = 65:	3CAC6514	969B2FC8	0D74807E	41C2091F	DAB08F2F
t = 66:	14CD9A35	3CAC6514	25A6CBF2	0D74807E	41C2091F
t = 67:	BA564047	14CD9A35	0F2B1945	25A6CBF2	0D74807E
t = 68:	C241F74D	BA564047	4533668D	0F2B1945	25A6CBF2
t = 69:	2896B70F	C241F74D	EE959011	4533668D	0F2B1945
t = 70:	564BBED1	2896B70F	70907DD3	EE959011	4533668D
t = 71:	8FA15D5A	564BBED1	CA25ADC3	70907DD3	EE959011
t = 72:	9A226C11	8FA15D5A	5592EFB4	CA25ADC3	70907DD3
t = 73:	F0B94489	9A226C11	A3E85756	5592EFB4	CA25ADC3
t = 74:	1809D5E2	F0B94489	66889B04	A3E85756	5592EFB4
t = 75:	B86C5A40	1809D5E2	7C2E5122	66889B04	A3E85756
t = 76:	DFE7E487	B86C5A40	86027578	7C2E5122	66889B04
t = 77:	70286C07	DFE7E487	2E1B1690	86027578	7C2E5122
t = 78:	24FF7ED5	70286C07	F7F9F921	2E1B1690	86027578
t = 79:	9A1F95A8	24FF7ED5	DC0A1B01	F7F9F921	2E1B1690.

After processing, the hex values of  $\{H_i\}$  are

$$H_0 = 67452301 + 9A1F95A8 = 0164B8A9$$

$$H_1 = EFCDAB89 + 24FF7ED5 = 14CD2A5E$$

$$H_2 = 98BADC9E + DC0A1B01 = 74C4F7FF$$

$$H_3 = 10325476 + F7F9F921 = 082C4D97$$

$$H_4 = C3D2E1F0 + 2E1B1690 = F1EDF880.$$

Message digest = 0164B8A9 14CD2A5E 74C4F7FF 082C4D97 F1EDF880

**APPENDIX B. A SECOND SAMPLE MESSAGE AND ITS MESSAGE DIGEST**

This appendix is for informational purposes only and is not required to meet the standard.

Let the message be the binary-coded form (cf. Appendix A) of the ASCII string

"abcdbcdecdefdefgefghfghighijhijkljklmklmnlnomnopnopq".

Since each of the 56 characters is converted to 8 bits, the length of the message is  $l = 448$ . In step (a) of Section 4, we append "1". In step (b) we append 511 "0"s. In step (c) we append the 2-word representation of 448, i.e., hex 00000000 000001C0. This gives  $n = 2$ .

The initial hex values of  $\{H_i\}$  are

$$H_0 = 67452301$$

$$H_1 = EFCDA89$$

$$H_2 = 98BADCCE$$

$$H_3 = 10325476$$

$$H_4 = C3D2E1F0.$$

Start processing block 1. The words of block 1 are

$W[0]$	= 61626364
$W[1]$	= 62636465
$W[2]$	= 63646566
$W[3]$	= 64656667
$W[4]$	= 65666768
$W[5]$	= 66676869
$W[6]$	= 6768696A
$W[7]$	= 68696A6B
$W[8]$	= 696A6B6C
$W[9]$	= 6A6B6C6D
$W[10]$	= 6B6C6D6E
$W[11]$	= 6C6D6E6F
$W[12]$	= 6D6E6F70
$W[13]$	= 6E6F7071
$W[14]$	= 80000000
$W[15]$	= 00000000.

The hex values of A,B,C,D,E after pass t of the "for t = 0 to 79" loop (step (d) of Section 7 or (c) of Section 8) are

	A	B	C	D	E
t = 0:	0116FC17	67452301	7BF36AE2	98BADC <sup>F</sup> E	10325476
t = 1:	EBF3B452	0116FC17	59D148C0	7BF36AE2	98BADC <sup>F</sup> E
t = 2:	5109913A	EBF3B452	C045BF05	59D148C0	7BF36AE2
t = 3:	2C4F6EAC	5109913A	BAFCED14	C045BF05	59D148C0
t = 4:	33F4AE5B	2C4F6EAC	9442644E	BAFCED14	C045BF05
t = 5:	96B85189	33F4AE5B	0B13DBAB	9442644E	BAFCED14
t = 6:	DB04CB58	96B85189	CCFD2B96	0B13DBAB	9442644E
t = 7:	45833F0F	DB04CB58	65AE1462	CCFD2B96	0B13DBAB
t = 8:	C565C35E	45833F0F	36C132D6	65AE1462	CCFD2B96
t = 9:	6350AFDA	C565C35E	D160CFC3	36C132D6	65AE1462
t = 10:	8993EA77	6350AFDA	B15970D7	D160CFC3	36C132D6
t = 11:	E19ECAA2	8993EA77	98D42BF6	B15970D7	D160CFC3
t = 12:	8603481E	E19ECAA2	E264FA9D	98D42BF6	B15970D7
t = 13:	32F94A85	8603481E	B867B2A8	E264FA9D	98D42BF6
t = 14:	B2E7A8BE	32F94A85	A180D207	B867B2A8	E264FA9D
t = 15:	42637E39	B2E7A8BE	4CBE52A1	A180D207	B867B2A8
t = 16:	66036329	42637E39	ACB9EA2F	4CBE52A1	A180D207
t = 17:	B59A89E4	66036329	5098DF8E	ACB9EA2F	4CBE52A1
t = 18:	90B9433E	B59A89E4	5980D8CA	5098DF8E	ACB9EA2F
t = 19:	DB5227E2	90B9433E	2D66A279	5980D8CA	5098DF8E
t = 20:	91241034	DB5227E2	A42E50CF	2D66A279	5980D8CA
t = 21:	4C06BD64	91241034	B6D489F8	A42E50CF	2D66A279
t = 22:	8665831E	4C06BD64	2449040D	B6D489F8	A42E50CF
t = 23:	3F62D9EC	8665831E	1301AF59	2449040D	B6D489F8
t = 24:	CD40E178	3F62D9EC	A19960C7	1301AF59	2449040D
t = 25:	D83E484E	CD40E178	0FD8B67B	A19960C7	1301AF59
t = 26:	D70940FE	D83E484E	3350385E	0FD8B67B	A19960C7
t = 27:	39B6981B	D70940FE	B60F9213	3350385E	0FD8B67B
t = 28:	694303AE	39B6981B	B5C2503F	B60F9213	3350385E
t = 29:	8E08FD0A	694303AE	CE6DA606	B5C2503F	B60F9213
t = 30:	FBFF1BA5	8E08FD0A	9A50C0EB	CE6DA606	B5C2503F
t = 31:	8AB96092	FBFF1BA5	A3823F42	9A50C0EB	CE6DA606
t = 32:	4206057A	8AB96092	7EFFC6E9	A3823F42	9A50C0EB
t = 33:	2CBCFC1A	4206057A	A2AE5824	7EFFC6E9	A3823F42
t = 34:	505759F3	2CBCFC1A	9081815E	A2AE5824	7EFFC6E9
t = 35:	05BB8EC9	505759F3	8B2F3F06	9081815E	A2AE5824
t = 36:	A0FC08A0	05BB8EC9	D415D67C	8B2F3F06	9081815E
t = 37:	8664F5E1	A0FC08A0	416EE3B2	D415D67C	8B2F3F06
t = 38:	FE3D2A4F	8664F5E1	283F0228	416EE3B2	D415D67C
t = 39:	07D02AA9	FE3D2A4F	61993D78	283F0228	416EE3B2
t = 40:	38D7321C	07D02AA9	FF8F4A93	61993D78	283F0228
t = 41:	1F3CA4C0	38D7321C	41F40AAA	FF8F4A93	61993D78
t = 42:	DF27AA0C	1F3CA4C0	0E35CC87	41F40AAA	FF8F4A93
t = 43:	84E2DBA6	DF27AA0C	07CF2930	0E35CC87	41F40AAA
t = 44:	8797EB77	84E2DBA6	37C9EA83	07CF2930	0E35CC87
t = 45:	9D220100	8797EB77	A138B6E9	37C9EA83	07CF2930
t = 46:	CB326B71	9D220100	E1E5FADD	A138B6E9	37C9EA83
t = 47:	505DE66F	CB326B71	27488040	E1E5FADD	A138B6E9
t = 48:	FFDF8E6F	505DE66F	72CC9ADC	27488040	E1E5FADD

t = 49:	47A17A6F	FFDF8E6F	D417799B	72CC9ADC	27488040
t = 50:	2C742CF4	47A17A6F	FFF7E39B	D417799B	72CC9ADC
t = 51:	692C82F3	2C742CF4	D1E85E9B	FFF7E39B	D417799B
t = 52:	741A7AEB	692C82F3	0B1D0B3D	D1E85E9B	FFF7E39B
t = 53:	E89625B3	741A7AEB	DA4B20BC	0B1D0B3D	D1E85E9B
t = 54:	BB527C29	E89625B3	DD069EBA	DA4B20BC	0B1D0B3D
t = 55:	609A8616	BB527C29	FA25896C	DD069EBA	DA4B20BC
t = 56:	5E259CED	609A8616	6ED49F0A	FA25896C	DD069EBA
t = 57:	FDCE04C4	5E259CED	9826A185	6ED49F0A	FA25896C
t = 58:	2A35958F	FDCE04C4	5789673B	9826A185	6ED49F0A
t = 59:	029A9DBB	2A35958F	3F738131	5789673B	9826A185
t = 60:	651604AB	029A9DBB	CA8D6563	3F738131	5789673B
t = 61:	3F163F73	651604AB	C0A6A76E	CA8D6563	3F738131
t = 62:	60E30527	3F163F73	D945812A	C0A6A76E	CA8D6563
t = 63:	DA53F35E	60E30527	CFC58FDC	D945812A	C0A6A76E
t = 64:	59F8E302	DA53F35E	/D838C149	CFC58FDC	D945812A
t = 65:	BE75732C	59F8E302	B694FCD7	D838C149	CFC58FDC
t = 66:	8D8DFD49	BE75732C	967E38C0	B694FCD7	D838C149
t = 67:	556247FC	8D8DFD49	2F9D5CCB	967E38C0	B694FCD7
t = 68:	C416C3E2	556247FC	63637F52	2F9D5CCB	967E38C0
t = 69:	64C244C9	C416C3E2	155891FF	63637F52	2F9D5CCB
t = 70:	B0DF5B97	64C244C9	B105B0F8	155891FF	63637F52
t = 71:	905723FE	B0DF5B97	59309132	B105B0F8	155891FF
t = 72:	49946022	905723FE	EC37D6E5	59309132	B105B0F8
t = 73:	B3A64DB3	49946022	A415C8FF	EC37D6E5	59309132
t = 74:	281589BC	B3A64DB3	92651808	A415C8FF	EC37D6E5
t = 75:	4623888D	281589BC	ECE9936C	92651808	A415C8FF
t = 76:	74EB04B7	4623888D	0A05626F	ECE9936C	92651808
t = 77:	035D4CD9	74EB04B7	5188E223	0A05626F	ECE9936C
t = 78:	B2BDD7D0	035D4CD9	DD3AC12D	5188E223	0A05626F
t = 79:	1D750196	B2BDD7D0	40D75336	DD3AC12D	5188E223.

Block 1 has been processed. The values of  $\{H_i\}$  are

$$H_0 = 67452301 + 1D750196 = 84BA2497$$

$$H_1 = EFCDAB89 + B2BDD7D0 = A28B8359$$

$$H_2 = 98BADCFE + 40D75336 = D9923034$$

$$H_3 = 10325476 + DD3AC12D = ED6D15A3$$

$$H_4 = C3D2E1F0 + 5188E223 = 155BC413.$$

Start processing block 2. The words of block 2 are

$$W[0] = 00000000$$

$$W[1] = 00000000$$

$$W[2] = 00000000$$

```

W[3] = 00000000
W[4] = 00000000
W[5] = 00000000
W[6] = 00000000
W[7] = 00000000
W[8] = 00000000
W[9] = 00000000
W[10] = 00000000
W[11] = 00000000
W[12] = 00000000
W[13] = 00000000
W[14] = 00000000
W[15] = 000001C0.

```

The hex values of A,B,C,D,E after pass t of the "for t = 0 to 79" loop (step (d) of Section 7 or (c) of Section 8) are

	A	B	C	D	E
t = 0:	D508E54E	84BA2497	68A2E0D6	D9923034	ED6D15A3
t = 1:	42AE69CC	D508E54E	E12E8925	68A2E0D6	D9923034
t = 2:	738C64E9	42AE69CC	B5423953	E12E8925	68A2E0D6
t = 3:	D5B4A0FE	738C64E9	10AB9A73	B5423953	E12E8925
t = 4:	870F3C0B	D5B4A0FE	5CE3193A	10AB9A73	B5423953
t = 5:	46574E97	870F3C0B	B56D283F	5CE3193A	10AB9A73
t = 6:	1405102F	46574E97	E1C3CF02	B56D283F	5CE3193A
t = 7:	297306DF	1405102F	D195D3A5	E1C3CF02	B56D283F
t = 8:	30185CE2	297306DF	C501440B	D195D3A5	E1C3CF02
t = 9:	10D7BA0C	30185CE2	CA5CC1B7	C501440B	D195D3A5
t = 10:	0C28CF6B	10D7BA0C	8C061738	CA5CC1B7	C501440B
t = 11:	6EABFEC0	0C28CF6B	0435EE83	8C061738	CA5CC1B7
t = 12:	7E85F170	6EABFEC0	C30A33DA	0435EE83	8C061738
t = 13:	F964F1A3	7E85F170	1BAAFFB0	C30A33DA	0435EE83
t = 14:	26E19055	F964F1A3	1FA17C5C	1BAAFFB0	C30A33DA
t = 15:	156937E7	26E19055	FE593C68	1FA17C5C	1BAAFFB0
t = 16:	6295F273	156937E7	49B86415	FE593C68	1FA17C5C
t = 17:	B81A706E	6295F273	C55A4DF9	49B86415	FE593C68
t = 18:	A5620A0D	B81A706E	D8A57C9C	C55A4DF9	49B86415
t = 19:	2DBC9CFF	A5620A0D	AE069C1B	D8A57C9C	C55A4DF9
t = 20:	BF89C409	2DBC9CFF	69588283	AE069C1B	D8A57C9C
t = 21:	239A6D9B	BF89C409	CB6F273F	69588283	AE069C1B
t = 22:	ADEC9CD5	239A6D9B	6FE27102	CB6F273F	69588283
t = 23:	1CDD463F	ADEC9CD5	C8E69B66	6FE27102	CB6F273F
t = 24:	E0DA5334	1CDD463F	6B7B2735	C8E69B66	6FE27102
t = 25:	B947BDAB	E0DA5334	C737518F	6B7B2735	C8E69B66
t = 26:	AD4E620C	B947BDAB	383694CD	C737518F	6B7B2735
t = 27:	CA67CF14	AD4E620C	EE51EF6A	383694CD	C737518F
t = 28:	FE343974	CA67CF14	2B539883	EE51EF6A	383694CD
t = 29:	7CFD680A	FE343974	3299F3C5	2B539883	EE51EF6A
t = 30:	E4D7304C	7CFD680A	3F8D0E5D	3299F3C5	2B539883

t = 31:	A6FD2352	E4D7304C	9F3F5A02	3F8D0E5D	3299F3C5
t = 32:	C57DADCD	A6FD2352	3935CC13	9F3F5A02	3F8D0E5D
t = 33:	5F146AB9	C57DADCD	A9BF48D4	3935CC13	9F3F5A02
t = 34:	469DC798	5F146AB9	715F6B73	A9BF48D4	3935CC13
t = 35:	03BCF3DA	469DC798	57C51AAE	715F6B73	A9BF48D4
t = 36:	F03F67BA	03BCF3DA	11A771E6	57C51AAE	715F6B73
t = 37:	2E04E8C4	F03F67BA	80EF3CF6	11A771E6	57C51AAE
t = 38:	E8B3497E	2E04E8C4	BC0FD9EE	80EF3CF6	11A771E6
t = 39:	A9CE9B40	E8B3497E	0B813A31	BC0FD9EE	80EF3CF6
t = 40:	F261BB65	A9CE9B40	BA2CD25F	0B813A31	BC0FD9EE
t = 41:	42EF9DD9	F261BB65	2A73A6D0	BA2CD25F	0B813A31
t = 42:	B2F2664A	42EF9DD9	7C986ED9	2A73A6D0	BA2CD25F
t = 43:	1291092A	B2F2664A	50BBE776	7C986ED9	2A73A6D0
t = 44:	7C6AEF48	1291092A	ACBC9992	50BBE776	7C986ED9
t = 45:	A9CB9DF6	7C6AEF48	84A4424A	ACBC9992	50BBE776
t = 46:	C5F82E71	A9CB9DF6	1F1ABBD2	84A4424A	ACBC9992
t = 47:	8868C238	C5F82E71	AA72E77D	1F1ABBD2	84A4424A
t = 48:	B052F768	8868C238	717E0B9C	AA72E77D	1F1ABBD2
t = 49:	61102AC0	B052F768	221A308E	717E0B9C	AA72E77D
t = 50:	8BEE2FF1	61102AC0	2C14BDDA	221A308E	717E0B9C
t = 51:	9E700133	8BEE2FF1	18440AB0	2C14BDDA	221A308E
t = 52:	877A43CD	9E700133	62FB8BFC	18440AB0	2C14BDDA
t = 53:	C4E901D6	877A43CD	E79C004C	62FB8BFC	18440AB0
t = 54:	2C7A07F0	C4E901D6	61DE90F3	E79C004C	62FB8BFC
t = 55:	67344973	2C7A07F0	B13A4075	61DE90F3	E79C004C
t = 56:	7EBAEE45	67344973	0B1E81FC	B13A4075	61DE90F3
t = 57:	EB9659B3	7EBAEE45	D9CD125C	0B1E81FC	B13A4075
t = 58:	0EBFB62A	EB9659B3	5FAEBB91	D9CD125C	0B1E81FC
t = 59:	4DBF216A	0EBFB62A	FAE5966C	5FAEBB91	D9CD125C
t = 60:	08089F12	4DBF216A	83AFED8A	FAE5966C	5FAEBB91
t = 61:	601ABA34	08089F12	936FC85A	83AFED8A	FAE5966C
t = 62:	E1685B50	601ABA34	820227C4	936FC85A	83AFED8A
t = 63:	EC956F26	E1685B50	1806AE8D	820227C4	936FC85A
t = 64:	6BED4126	EC956F26	385A16D4	1806AE8D	820227C4
t = 65:	96D6E5E6	6BED4126	BB255BC9	385A16D4	1806AE8D
t = 66:	A5D83970	96D6E5E6	9AFB5049	BB255BC9	385A16D4
t = 67:	74CCF6E4	A5D83970	A5B5B979	9AFB5049	BB255BC9
t = 68:	B9BDCA6D	74CCF6E4	29760E5C	A5B5B979	9AFB5049
t = 69:	9526A197	B9BDCA6D	1D333DB9	29760E5C	A5B5B979
t = 70:	A2E5A7C9	9526A197	6E6F729B	1D333DB9	29760E5C
t = 71:	3708B81B	A2E5A7C9	E549A865	6E6F729B	1D333DB9
t = 72:	F27081EC	3708B81B	68B969F2	E549A865	6E6F729B
t = 73:	41DAEB9B	F27081EC	CDC22E06	68B969F2	E549A865
t = 74:	4215A57B	41DAEB9B	3C9C207B	CDC22E06	68B969F2
t = 75:	2655C2D6	4215A57B	D076BAE6	3C9C207B	CDC22E06
t = 76:	11DC8A86	2655C2D6	D085695E	D076BAE6	3C9C207B
t = 77:	69364641	11DC8A86	899570B5	D085695E	D076BAE6
t = 78:	0A6ED856	69364641	847722A1	899570B5	D085695E
t = 79:	4D974A4A	0A6ED856	5A4D9190	847722A1	899570B5

## APPENDIX C: A THIRD SAMPLE MESSAGE AND ITS MESSAGE DIGEST

Block 2 has been processed. The values of  $\{H_i\}$  are

$$H_0 = 84BA2497 + 4D974A4A = D2516EE1$$

$$H_1 = A28B8359 + 0A6ED856 = ACFA5BAF$$

$$H_2 = D9923034 + 5A4D9190 = 33DFC1C4$$

$$H_3 = ED6D15A3 + 847722A1 = 71E43844$$

$$H_4 = 155BC413 + 899570B5 = 9EF134C8.$$

Message digest = D2516EE1 ACFA5BAF 33DFC1C4 71E43844 9EF134C8

**APPENDIX C. A THIRD SAMPLE MESSAGE AND ITS MESSAGE DIGEST**

This appendix is for informational purposes only and is not required to meet the standard.

Let the message be the binary-coded form of the ASCII string which consists of 1,000,000 repetitions of "a".

**Message digest = 3232AFFA 48628A26 653B5AAA 44541FD9 0D690603**

a = 01: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
b = 1: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
c = 2: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
d = 3: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
e = 4: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
f = 5: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
g = 6: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
h = 7: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
i = 8: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
j = 9: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
k = 10: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
l = 11: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
m = 12: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
n = 13: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
o = 14: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
p = 15: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
q = 16: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
r = 17: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
s = 18: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
t = 19: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
u = 20: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
v = 21: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
w = 22: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
x = 23: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
y = 24: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
z = 25: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
0 = 26: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
1 = 27: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
2 = 28: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
3 = 29: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
4 = 30: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
5 = 31: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
6 = 32: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
7 = 33: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
8 = 34: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
9 = 35: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
0 = 36: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
1 = 37: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
2 = 38: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
3 = 39: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
4 = 40: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
5 = 41: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
6 = 42: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
7 = 43: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
8 = 44: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
9 = 45: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
0 = 46: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
1 = 47: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
2 = 48: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
3 = 49: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
4 = 50: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
5 = 51: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
6 = 52: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
7 = 53: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
8 = 54: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
9 = 55: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
0 = 56: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
1 = 57: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
2 = 58: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
3 = 59: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
4 = 60: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
5 = 61: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
6 = 62: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
7 = 63: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
8 = 64: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
9 = 65: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
0 = 66: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
1 = 67: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
2 = 68: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
3 = 69: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
4 = 70: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
5 = 71: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
6 = 72: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
7 = 73: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
8 = 74: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
9 = 75: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
0 = 76: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
1 = 77: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
2 = 78: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603
3 = 79: 3232AFFA	48628A26	653B5AAA	44541FD9	0D690603

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